



## The basis number of connected vertex-disjoint graphs

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### Abstract

The basis number  $b(G)$  of a graph  $G$  is defined to be the smallest positive integer  $k$  such that  $G$  has a  $k$ -fold basis for its cycle space. We try to find an upper bound for  $b(G_1 + G_2 + G_3 + G_4)$ .

We prove that, if  $G_1, G_2, G_3$  and  $G_4$  are connected vertex-disjoint graphs and each has a spanning tree of vertex degree not more than 4, then  $b(G_1 + G_2 + G_3 + G_4) \leq \max\{4, b(G_1) + 1, b(G_2) + 2, b(G_3) + 2, b(G_4) + 1\}$ .

The basis number of quadruple join of paths will be studied, where we prove that

$$b(p_m + p_n + p_p + p_t) = 4, \quad \forall m, t \geq 5 \text{ and } n, p \geq 6.$$

**Keywords:** Graph; Basis number; Connected vertex-disjoint graphs; Path

### 1. Introduction

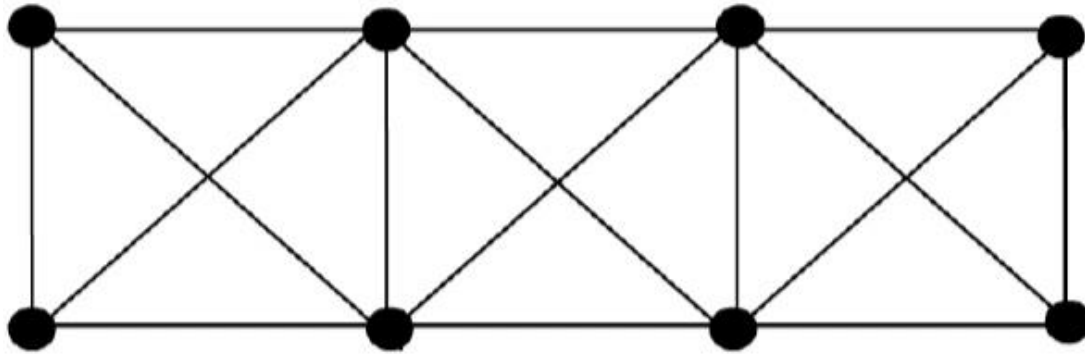
In the literature of graph theory, the interest in the basic number of data has increased; we refer the reader back to the research [4-11]. In this research, we will assume that all the data that we encounter are finite, undirected, and simple data; for terms that are not defined (see [9] and [10]). Let  $G$  be a statement connected by its edges  $e_1, e_2, e_3, \dots, e_q$ . For every partial  $s$  set of edges of  $G$ , there is a vector  $(a_1, a_2, a_3, \dots, a_q)$  corresponds to  $S$  such that  $a_i = 1$  if  $e_i \in S$ . These vectors are a vector space of dimension  $q$  on the field  $\mathbf{Z}_2$ , it is called the vector space associated with the statement  $G$  and is denoted by  $[(\mathbf{Z}_2)^q]$ . The vectors  $(\mathbf{Z}_2)^q$  corresponding to the circuits of  $G$  generate a partial vector space called **the circuit space** of  $G$ , denoted by  $\xi(G)$ . Each vector in  $\xi(G)$  represents either a circuit in  $G$  or a union of circuits separated from each other relative to the edges. One of the well-known results in the statement theorem is that the dimension of the circuit space of a statement connected to  $\xi(G)G$  is  $q-p+1$ , since  $p$  represents the number of vertices of  $G$  and  $Q$  the number of its edges. The way to find the base for the space of circuits  $\xi(G)$  is as follows:

- Let  $T$  be a generator tree in  $G$ ; If  $e$  is an edge belonging to  $G-T$  then  $T+e$  contains only one circuit and let be  $C_e$ . That  $q-p+1$  of the circuits  $C_e$ , where  $e \in G - T$  forms the base of the circuit space  $\xi(G)$ .
- A base  $B$  in the space of circuits  $\xi(G)$  is said to be  **$k$ -fold** if each edge of  $G$  does not appear more than  $k$  times in the circuits corresponding to the vectors in base  $B$ .
- The base number of the statement  $G$  is defined to be the smallest integer  $k$  such that  $\xi(G)$  has a base with a  $k$ -fold; and is denoted by  $b(G)$ . If  $B$  is a base for the space of circuits  $\xi(G)$  and  $e$  is an edge of  $G$ , then **the fold of the edge  $e$**  in  $B$  is denoted by  $f_B(e)$  is the number of circuits in  $B$  containing edge  $e$ .

### 2. Definition (1)

Let  $G_1 = (X, E_1)$ ,  $G_2 = (Y, E_2)$ ,  $G_3 = (Z, E_3)$ ,  $G_4 = (U, E_4)$  simple data and separate for warheads, the data connection  $G_1$  and  $G_2$  is a data set of its vertices is  $X \cup Y$  and the group of the IS  $E_1 \cup E_2 \cup \{xy | x \in X, y \in Y\}$  is denoted by  $G_1 + G_2$ ; as well as **the tripartite contact** data  $G_1, G_2, G_3$  is a data set of its vertices is  $X \cup Y \cup Z$  And the sum of its edges is  $E_1 \cup E_2 \cup E_3 \cup \{xy | x \in X, y \in Y\} \cup \{yz | y \in Y, z \in Z\}$  is denoted by  $G_1 + G_2 + G_3$ ; similarly, we call **the**

**quadrilateral contact**  $G_1 + G_2 + G_3 + G_4$  that the data set of its vertices is  $X \cup Y \cup Z \cup U$  and the sum of its edges is  $E_1 \cup E_2 \cup E_3 \cup E_4 \cup \{xy|x \in X, y \in Y\} \cup \{yz|y \in Y, z \in Z\} \cup \{zu|z \in Z, u \in U\}$   
 For example,  $P_4 = K_1 + K_1 + K_1 + K_1$  and Figure 1 shows the quadrilateral contact of the complete statement  $K_2$



**Figure 1.** The quadrilateral contact of the complete statement  $K_2$ .

Form (1).  $K_2 + K_2 + K_2 + K_2$  And in a observational way we know the connection n of the data,  $G_1 + G_2 + \dots + G_n$ ; it is clear that this statement is connected. We will encode the path that has n vertices with the symbol  $P_n$   
 In 1937, McClain [13] proved the first important theorem on the subject, which states that "a statement G is planar if and only if  $b(G) \leq 2$ ". In 1981 and 1982, Schleicher ([8], [16]) found the base number of the perfect statement  $K_n$ , the perfect binary hash statement  $K_{m,n}$  and a cube-n. The base number of the lexical novelty of the data was calculated in [1], as well as the base number of the Cartesian novelty of some data was studied in [2], and Ali studied [3] the base number of the connection of two data where he proved if  $G_1$  and  $G_2$  are two separate data for vertices and each of them has a generating forest with a degree not exceeding 4, then:

$$b(G_1 + G_2) \leq \max\{4, b(G_1) + 1, b(G_2) + 1\}.$$

In addition, in [15] the basic number of the tripartite contact of data was studied. In this research, we will study the fundamental number of the quadratic contact of data, where we proved:

$$b(G_1 + G_2 + G_3 + G_4) \leq \max\{4, b(G_1) + 1, b(G_2) + 2, b(G_3) + 2, b(G_4) + 1\}$$

And

$$b(P_m + P_n + P_p + P_t) = 4, \quad \forall m, t \geq 5 \quad \text{and} \quad n, p \geq 6$$

**The base number of  $G_1 + G_2 + G_3 + G_4$ :**

Suppose that  $X = \{x_1, x_2, \dots, x_m\}$ ,  $Y = \{y_1, y_2, \dots, y_n\}$ ,  $Z = \{z_1, z_2, \dots, z_p\}$ ,  $U = \{u_1, u_2, \dots, u_t\}$  we suppose that  $K_{m,n}$ ,  $K_{n,p}$ ,  $K_{p,t}$  are the binary hash data whose independent sets, respectively, are  $Y, X$ ,  $Y, Z$ ,  $Z, U$ . Schmeichel [16] proved that:

$$B_r(K_{m,n}) = \{x_i y_j x_{i+1} y_{j+1} | i = 1, 2, \dots, m-1 \text{ and } j = 1, 2, \dots, n-1\}, \quad ..(1)$$

$$B_r(K_{n,p}) = \{y_i z_j y_{i+1} z_{j+1} | i = 1, 2, \dots, n-1 \text{ and } j = 1, 2, \dots, p-1\}, \quad ..(2)$$

$$B_r(K_{p,t}) = \{z_i u_j z_{i+1} u_{j+1} | i = 1, 2, \dots, p-1 \text{ and } j = 1, 2, \dots, t-1\}, \quad ..(3)$$

Are the required bases for  $\xi(K_{m,n})$ ,  $\xi(K_{n,p})$ ,  $\xi(K_{p,t})$  with 4 folds respectively; that is,  $b(K_{m,n}) \leq 4$ ,  $b(K_{n,p}) \leq 4$ ,  $b(K_{p,t}) \leq 4$  for each  $m, n, p, t \geq 5$ ; Also, Ali [3] proved that:

$$A_1 = \{x_1 y_1, x_1 y_n, x_m y_1, x_m y_n\} \cup \{y_1 z_1, y_1 z_p, y_n z_1, y_n z_p\} \cup \{z_1 u_1, z_1 u_t, z_p u_1, z_p u_t\} \quad ..(4)$$

Are the sets of edges  $K_{m,n}$ ,  $K_{n,p}$ ,  $K_{p,t}$  with a 1-fold in the required base  $B_r(K_{m,n})$ ,  $B_r(K_{n,p})$ ,  $B_r(K_{p,t})$  respectively and that:

$$A_2 = \left( \{x_1 y_j, x_m y_j | j = 2, 3, \dots, n-1\} \cup \{x_i y_1, x_i y_n | i = 2, 3, \dots, m-1\} \right) \cup \left( \{y_1 z_j, y_n z_j | j = 2, 3, \dots, p-1\} \cup \{y_j z_1, y_j z_p | j = 2, 3, \dots, n-1\} \right) \cup \left( \{z_1 u_j, z_p u_j | j = 2, 3, \dots, t-1\} \cup \{z_j u_1, z_j u_t | j = 2, 3, \dots, p-1\} \right) \quad ..(5)$$

Is the set of edges  $K_{m,n}$ ,  $K_{n,p}$ ,  $K_{p,t}$  with a fold of 2 in the required base  $B_r(K_{m,n})$ ,  $B_r(K_{n,p})$ ,  $B_r(K_{p,t})$  respectively, there is no edge with a fold of 3 in  $B_r$  and the rest of the edges are bent 4.

Now we will study the base number of the quadratic contact of the data, that is,  $b(G_1 + G_2 + G_3 + G_4)$  when  $G_1, G_2, G_3, G_4$  are connected statements and have a generating tree with a degree of not more than 4.

Let  $T_1, T_2, T_3, T_4$  be generating trees for  $G_1, G_2, G_3, G_4$  in such an order that for  $i = 1, 2, 3, 4$ , then:

$$deg_{T_i} u \leq 4, \forall u \in V(T_i)$$

Since every tree has at least two vertices with degree one, we assume that:

$$deg_{T_1}(x_1) = deg_{T_1}(x_m) = deg_{T_2}(y_1) = deg_{T_2}(y_n) = deg_{T_3}(z_1) = deg_{T_3}(z_p) = deg_{T_4}(u_1) = deg_{T_4}(u_t) = 1$$

Now for each edge  $x_i x_j$  in the tree  $T_1$  we will choose one triple circuit and only one by the formula  $x_i x_j y_1$  or  $x_i x_j y_n$  and this is in the following way:

- We divide the set of edges of  $T_1$  into two partial sets,  $W_1$  and  $W_1'$ , so that each vertex of  $x$  in  $T_1$  is the number of edges located on it that belong to the same partial set of not more than 2.

Since  $deg_{T_1}(x) \leq 4$  for every  $x$  in  $T_1$ , this hash is possible. Now, let  $S_1$  be the set of three circuits of  $T_1$  defined in the following form:

$$S_1 = \{x_i x_j y | x_i x_j \in E(T_1), y = y_i \text{ if } x_i x_j \in W_1 \text{ and } y = y_n \text{ if } x_i x_j \in W_1'\}$$

In the same way, we obtain the sets  $S_2, S_3, S_4$  from triangular circuits of tree edges  $T_2, T_3, T_4$  from the data  $G_2, G_3, G_4$ , respectively, where:

$$S_2 = \{x_i y_j y, y_i y_j z | y_i y_j \in E(T_2), x = x_1, z = z_1 \text{ if } y_i y_j \in W_2 \text{ and } x = x_m, z = z_p \text{ if } y_i y_j \in W_2'\}$$

$$S_3 = \{y_i z_j z, z_i z_j u | z_i z_j \in E(T_3), y = y_1, u = u_1 \text{ if } z_i z_j \in W_3 \text{ and } y = y_n, u = u_t \text{ if } z_i z_j \in W_3'\} \text{ and}$$

$$S_4 = \{z u_i u_j | u_i u_j \in E(T_4), z = z_1 \text{ if } u_i u_j \in W_4 \text{ and } z = z_p \text{ if } u_i u_j \in W_4'\}$$

It is clear that  $|S_1| = m - 1, |S_2| = 2(n - 1), |S_3| = 2(p - 1), |S_4| = t - 1$  and that

Each of the edges  $x_i y_1, x_i y_n$  for the values  $i = 1, 2, \dots, m$  appears at most in two circles in  $S_1$ .

And each of the edges  $y_j x_1, y_j x_m$  for the values  $j = 1, 2, \dots, n$  appears at most in two circles in  $S_2$ .

And each of the edges  $y_j z_1, y_j z_p$  for the values  $j = 1, 2, \dots, n$  appears at most in two circles in  $S_2$ .

And each of the edges  $z_j y_1, z_j y_n$  for the values  $j = 1, 2, \dots, p$  appears at most in two circles in  $S_3$ .

And each of the edges  $z_j u_1, z_j u_t$  for the values  $j = 1, 2, \dots, p$  appears at most in two circles in  $S_3$ .

Finally:

Each of the edges  $u_j z_1, u_j z_p$  for the values  $j = 1, 2, \dots, t$  appear at most in two circles in  $S_4$

And now, we are ready to prove the basic theorem in this research.

### 3. Theorem (1)

If  $G_1, G_2, G_3, G_4$  are connected and separate data for the vertices and each of them has a generating tree with a degree of not more than 4, then:

$$b(G_1 + G_2 + G_3 + G_4) \leq \max\{4, b(G_1) + 1, b(G_2) + 2, b(G_3) + 2, b(G_4) + 1\}.$$

**Proof:**

Let  $B_r(G_1), B_r(G_2), B_r(G_3), B_r(G_4)$  be the required bases for the data  $G_1, G_2, G_3, G_4$ , respectively; in addition, let  $B_r(K_{m,n}), B_r(K_{n,p}), B_r(K_{p,t})$  be the required bases for the data  $K_{m,n}, K_{n,p}, K_{p,t}$  are defined in (1), (2), (3) respectively.

Let:

$$B = B_r(G_1) \cup B_r(G_2) \cup B_r(G_3) \cup B_r(G_4) \cup B_r(K_{m,n}) \cup B_r(K_{n,p}) \cup B_r(K_{p,t}) \cup S_1 \cup S_2 \cup S_3 \cup S_4$$

Now we will prove that B is a norm of the space of circuits  $\xi(G_1 + G_2 + G_3 + G_4)$ ; it is obvious that:

$$\begin{aligned} |B| &= \dim \xi(G_1) + \dim \xi(G_2) + \dim \xi(G_3) + \dim \xi(G_4) + \dim \xi(K_{m,n}) + \dim \xi(K_{n,p}) + \dim \xi(K_{p,t}) + (m - 1) + \\ &2(n - 1) + 2(p - 1) + (t - 1) = (q_1 - m + 1) + (q_2 - n + 1) + (q_3 - p + 1) + (q_4 - t + 1) + mn - (m + \\ &n) + 1 + np - (n + p) + 1 + pt - (p + t) + 1 + m - 1 + 2n - 2 + 2p - 2 + t - 1 = (q_1 + q_2 + q_3 + q_4 + \\ &mn + np + pt) - (m + n + p + t) + 1 \\ &= \dim \xi(G_1 + G_2 + G_3 + G_4) \end{aligned}$$

Where  $q_1, q_2, q_3, q_4$  are the number of edges in  $G_1, G_2, G_3, G_4$  respectively.

We will prove that B is a linearly independent set of circuits; it is obvious that:

$$B_r(G_1) \cup B_r(G_2) \cup B_r(G_3) \cup B_r(G_4) \cup B_r(K_{m,n}) \cup B_r(K_{n,p}) \cup B_r(K_{p,t})$$

A linearly independent set and that both sets  $S_1, S_2, S_3, S_4$  are independent because they represent the boundaries of regions in a planar statement; now if C is any circuit generated by circuits  $S_1$ , C has an edge of  $T_1$ ; this edge does not

exist in any linear combination of circuits  $S_2$ . Therefore,  $S_1 \cup S_2$  independent; as well as circuits  $S_3$  be independent of the  $S_1 \cup S_2$  because each circuit of  $S_3$  contain on the edge of the  $T_3$ ; this edge is not present in any circuit of circuits  $S_1 \cup S_2$  therefore  $S_1 \cup S_2 \cup S_3$  circuits are linearly independent; and so the circuits  $S_4$  independent circuits  $S_1 \cup S_2 \cup S_3$  because any circuit in  $S_4$  contain on the edge of  $T_4$ ; this edge is not present in any linear structure of  $S_1 \cup S_2 \cup S_3$  thus be  $S_1 \cup S_2 \cup S_3 \cup S_4$  circuits are linearly independent. And for the proof of the set  $S = S_1 \cup S_2 \cup S_3 \cup S_4$  is linearly independent of the set.

$$B' = B_r(G_1) \cup B_r(G_2) \cup B_r(G_3) \cup B_r(G_4) \cup B_r(K_{m,n}) \cup B_r(K_{n,p}) \cup B_r(K_{p,t})$$

Note that any circuit generated from S contains edges from  $G_i$  for the values of  $i=1,2,3,4$  and also contains edges from  $K_{j,k}$  for the values of  $j=m, n, p$  and  $k = n, p, t$ ; in contrast, any circuit generated from B' contains edges from  $G_i$  for the values of  $i =1,2,3,4$  or edges from  $K_{j,k}$  for the values of  $j=m, n, p$  and  $k = n, p, t$ ; because it is a discrete set for edges, that is:

$$\sum_{i=1}^{|S|} C_i \neq \sum_{j=1}^{|B'|} b_j, C_i \in S \text{ and } b_j \in B'$$

Thus,  $B = S \cup B'$  is a linearly independent set and therefore a base for  $\xi(G_1 + G_2 + G_3 + G_4)$ .

Now we calculate the Tuck for each of the edges of the statement; from (5), The Tuck for each of the edges below:

$$\begin{aligned} x_i y_1, x_i y_n & , i = 2, 3, \dots, m - 1, \\ x_1 y_j, x_m y_j & , j = 2, 3, \dots, n - 1, \\ y_1 z_j, y_n z_j & , j = 2, 3, \dots, p - 1, \\ y_j z_1, y_j z_p & , j = 2, 3, \dots, n - 1, \\ z_1 u_j, z_p u_j & , j = 2, 3, \dots, t - 1, \\ z_j u_1, z_j u_t & , j = 2, 3, \dots, p - 1. \end{aligned}$$

In the next  $B_r(K_{m,n}) \cup B_r(K_{n,p}) \cup B_r(K_{p,t})$  is 2 and each edge appears at most at two of  $S_1 \cup S_2 \cup S_3 \cup S_4$ . Also, from (4) tuck to each of the edges below:

$$x_1 y_1, x_1 y_n, x_m y_1, x_m y_n, y_1 z_1, y_1 z_p, y_n z_1, y_n z_p, z_1 u_1, z_1 u_t, z_p u_1, z_p u_t$$

Is 1 in  $B_r(K_{m,n}) \cup B_r(K_{n,p}) \cup B_r(K_{p,t})$  each edge appears at most at two of  $S_1 \cup S_2 \cup S_3 \cup S_4$ ; therefore:

$$f_B(e) \leq 4, \forall e \in K_{m,n} \cup K_{n,p} \cup K_{p,t}$$

In addition, it is easy to verify that:

$$f_B(e) \leq \begin{cases} b(G_1) + 1, & \forall e \in G_1, \\ b(G_2) + 2, & \forall e \in G_2, \\ b(G_3) + 2, & \forall e \in G_3, \\ b(G_4) + 1, & \forall e \in G_4, \end{cases}$$

So, the Tuck of base B is  $\max\{4, b(G_1) + 1, b(G_2) + 2, b(G_3) + 2, b(G_4) + 1\}$   
Thus, the proof is done.

#### 4. Result

For every positive integer  $m, n, p, t \geq 3, b(P_m + P_n + P_p + P_t) \leq 4$ , equality is achieved when  $m, t \geq 5$  and  $n, p \geq 6$ .

##### Proof:

From the theorem (1):

$$b(P_m + P_n + P_p + P_t) \leq 4 \dots(6)$$

It is easy to prove that  $P_m + P_n + P_p + P_t$  is an uneven statement and therefore, according to McLean [13],  $b(P_m + P_n + P_p + P_t) \leq 4$  for the values  $m, n, p, t \geq 3$ .

Suppose that  $b(P_m + P_n + P_p + P_t) = 3$ ; let  $B_r$  be the required base for  $P_m + P_n + P_p + P_t$ , so:

$$|B_r| = |E(P_m + P_n + P_p + P_t)| - |V(P_m + P_n + P_p + P_t)| + 1$$

$$|B_r| = m - 1 + n - 1 + p - 1 + t - 1 + mn + np + pt - (m + n + p + t) + 1 = mn + np + pt - 3$$

Now the number of circuits of length 3 in the required base  $B_r$  is:

$$(m-1) + 2(n-1) + 2(p-1) + t-1 = m + 2n + 2p + t - 6$$

And the rest of the circuits with a length of 4 or more, since:

$$\sum_{i=1}^{dim\xi(G)} |C_i| \leq b \cdot |E(G)|$$

Where  $|C_i|$  stands for the length of the circuit  $C_i$  and  $b$  is the base number of the statement  $G$ , then:

$$4[(mn + np + pt - 3) - (m + 2n + 2p + t - 6)] + 3(m + 2n + 2p + t - 6) \leq 3[(m - 1) + (n - 1) + (p - 1) + (t - 1) + mn + np + pt]$$

That is;

$$mn + np + pt + 6 \leq 4m + 5n + 5p + 4t$$

Now, we can easily verify that the above inequality does not hold when  $m, t \geq 5$  and  $n, p \geq 6$ , so that:

$$b(p_m + p_n + p_p + p_t) = 4, \quad \forall m, t \geq 5 \text{ and } n, p \geq 6 \quad \dots(7)$$

In addition, by connecting (6) with (7) we get equality.

**Observation:**

In a similar way to the proof of theorem 1, we prove that if  $G_1, G_2, \dots, G_n$  is continuous and discrete data relative to the vertices and each of them has a generating tree with Degree Not exceeding 4, for each  $n \geq 3$

$$b(G_1 + G_2 + \dots + G_n) \leq \max\{4, b(G_1) + 1, b(G_n) + 1, b(G_i) + 2 | i = 2, 3, \dots, n - 1\}$$

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